

# UNIVERSITÄT

#### Searching for Genomic Variants in Multiple Genomes at Once

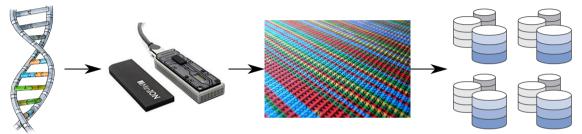
Jens Stove

Faculty of Technology, Center for Biotechnology (CeBiTec), Bielefeld Institute for Bioinformatics Infrastructure (BIBI) Bielefeld University. Germany

(joint work with Tizian Schulz, Roland Wittler, Sven Rahmann, Faraz Hach)

#### Motivation

#### Current status in DNA sequence analysis:





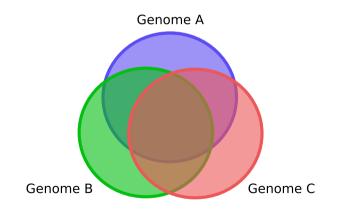
- Numerous individual genomes per species are sequenced
- $\rightarrow\,$  Huge amounts of partially redundant sequences

#### Challenges:

- Efficient storage
- Efficient analysis
- Effective analysis

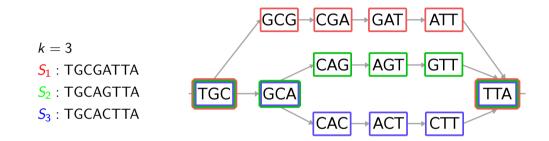
#### The sequence-based pangenome

How to efficiently store huge amounts of partially redundant sequences?



# Colored de Bruijn graphs

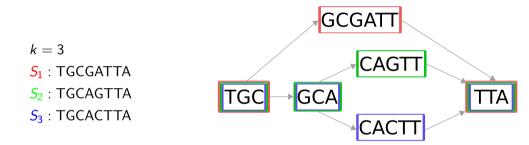
Colored de Bruijn graph (C-DBG)



- Vertices represent substrings of length k (k-mers)
- Associated color represents sequence of origin
- Edges between vertices that share a k-1 overlap

# Colored de Bruijn graphs

Compacted colored de Bruijn graph

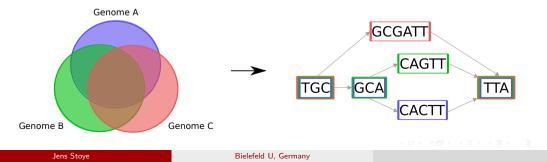


- Vertices of unique paths are replaced by single vertex (unitig)
- Label of unitig is sequence spelled by the path
- Allows even more efficiency

# Colored de Bruijn graphs (C-DBGs)

Why to store pan-genomes as compacted colored de Bruijn graphs?

- Reduced memory requirements
- No data preprocessing necessary (assembly)
- Data is stored based on similarity
- Allows analysis of multiple sequences in parallel



### Querying a sequence DB - BLAST

#### **BLAST** – Basic Local Alignment Search Tool

Developed by Stephen Altschul *et al.* in 1990 (second version in 1997) Is used to query databases of DNA and protein sequences

Different flavors are available:

- DNA→DNA (blastn)
- protein→protein (blastp)
- translated DNA→protein (blastx)
- protein→translated DNA (tblastn)

Algorithmic steps:

- Finding BLAST hits
- I Hit extension X-drop algorithm
- Gapped alignment calculation

....

### Querying a sequence DB - BLAST

#### **BLAST – Basic Local Alignment Search Tool**

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#### Problem statement:

Given a query sequence  $x \in \Sigma^*$  and a sequence database Y, find all highest scoring local alignments between x and  $y \in Y$  above a certain significance value and return them along with their score.

How to query a C-DBG to analyze the pangenome it represents?

**Idea:** Develop a method that queries a C-DBG in a BLAST-like manner to find local alignments between a query sequence and the sequences stored in the C-DBG.

 $\rightarrow$  Pangenome Local Alignment Search Tool – PLAST

#### **Problem statement:**

Given a query sequence  $x \in \Sigma^*$  and a set of sequences Y represented as a C-DBG, find all highest scoring local alignments between x and  $y \in Y$  above a certain significance value and return them along with their score.

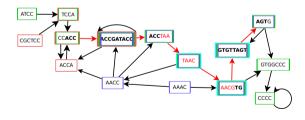
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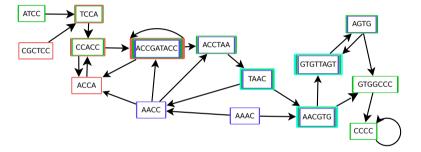
 $\rightarrow$  Pangenome Local Alignment Search Tool – PLAST

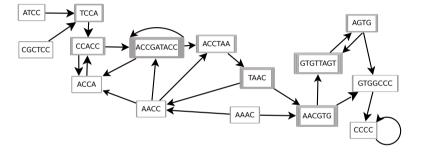
Procedure of PLAST:

- Seed detection
- Seed extension
- Gapped alignment calculation

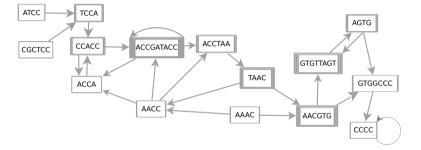


#### *x*: A G G A C C G - T A - - T A A C G G G G A A G T A C T *y*: A C C G A T A C C T A A C G T G T T A G T

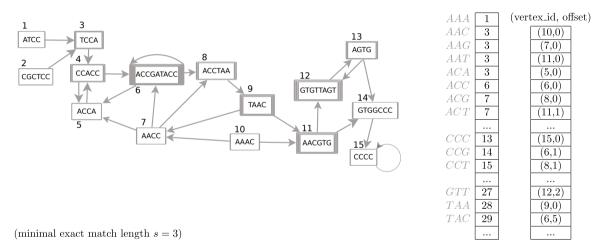




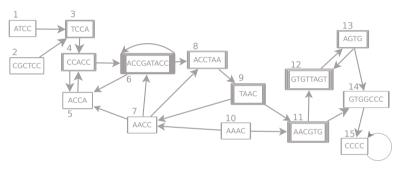
8



8



1. Seed detection:



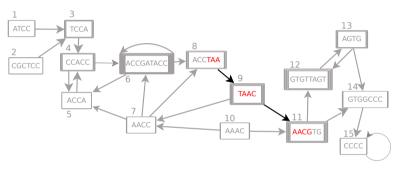


$(vertex\_id,offset)$	
(1	$^{0,0)}$
(7	(,0)
(1	$^{1,0)}$
(5	(0,0)
(6	5,0)
(8	$^{3,0)}$
(1	$^{1,1)}$
	$^{5,0)}$
	5,1)
(8	$^{3,1)}$
	$^{2,2)}$
	0,0)
(6	5,5)

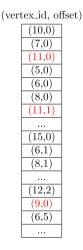
x: A G G A C C G T A T A A C G G G G A A G T A C T

(minimal exact match length s = 3)

1. Seed detection:



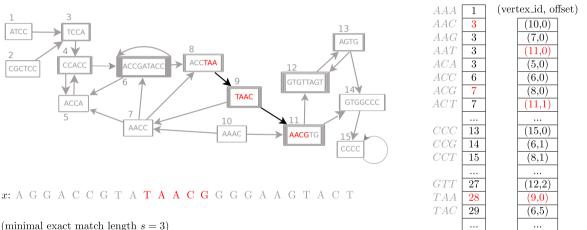




x: A G G A C C G T A **T A A C G** G G G A A G T A C T

(minimal exact match length s = 3)

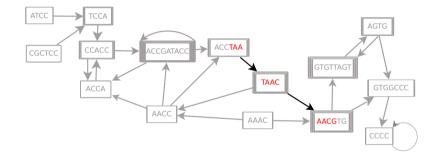
1. Seed detection:



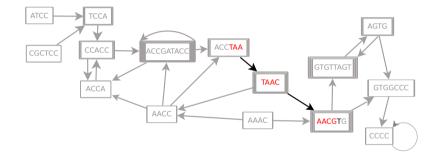
Possible in the same way BLAST does, using a precalculated index of the graph

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2. Seed extension:

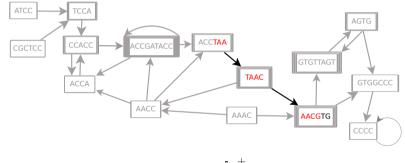


2. Seed extension:

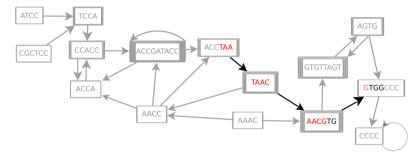




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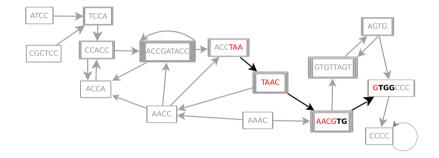


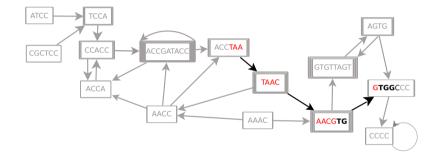






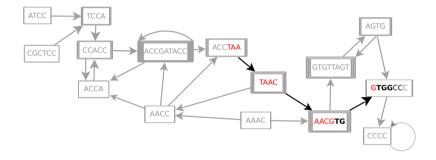
2. Seed extension:



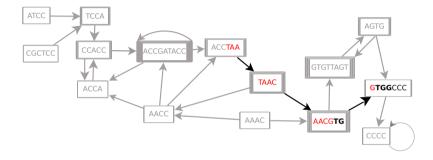




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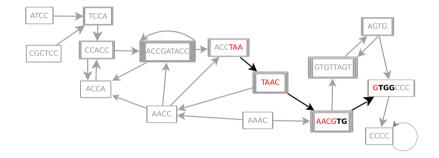


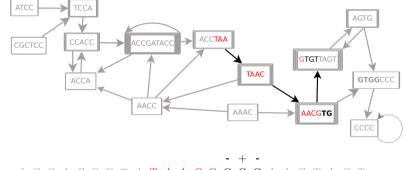






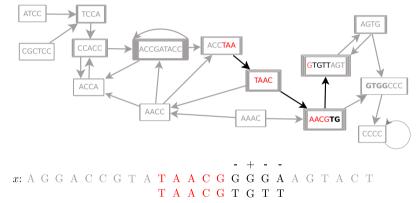
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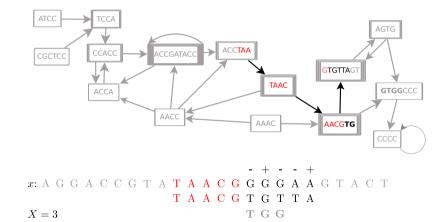


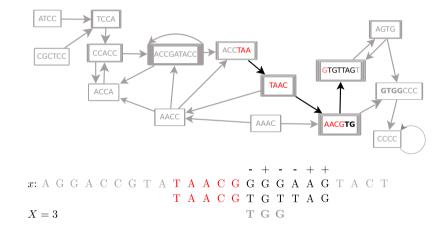


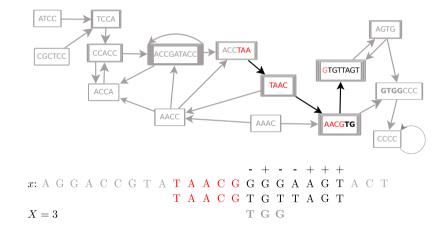
2. Seed extension:



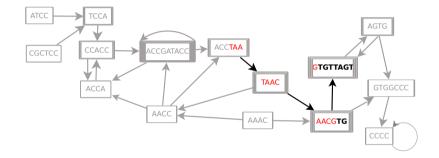
X = 3 T G G





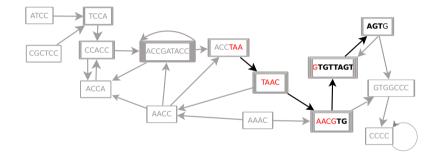


2. Seed extension:



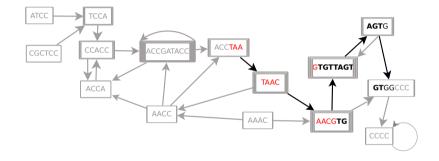


2. Seed extension:



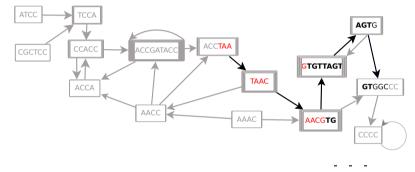
*x*: A G G A C C G T A **T A A C G G G G A A G T** A C T **T A A C G T G T T A G T** G

2. Seed extension:



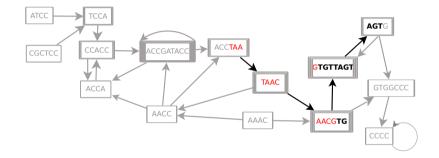


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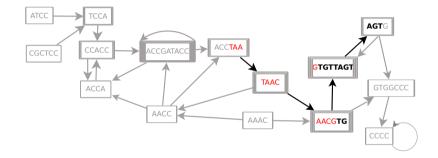


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*x*: A G G A C C G T A **T A A C G G G G A A G T** A C T **T A A C G T G T T A G T** 

X = 3Expensive since graph sequence is ambiguous

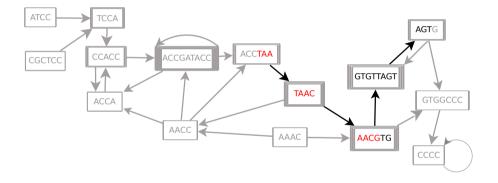
Brute-force traversal (DFS) mostly feasible

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#### The endless extension problem:

- Graphs may have some collapsed regions
- Regions have
  - many short loops
  - vertices with short sequences and high branching factor
- $\rightarrow$  nearly every possible sequence may be generated here
- $\rightarrow$  X-drop algorithm is unable to end extensions  $\Rightarrow$  hard break

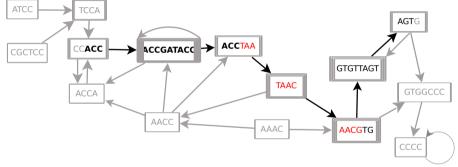
3. Gapped alignment of best extended seeds:



#### 

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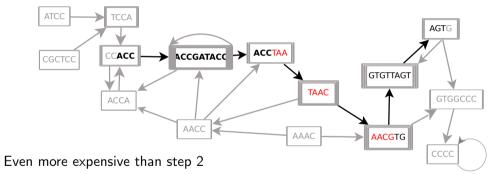
3. Gapped alignment of best extended seeds:



Banded, gapped alignment calculation following the extension path

```
x: A G G A C C G - T A - T A A C G G G A A G T A C T 
A C C G A T A C C T A A C G T G T T A G T
```

3. Gapped alignment of best extended seeds:



Only performed for a small subset of all findings, thus feasible

```
x: A G G A C C G - T A - T A A C G G G A A G T A C T 
A C C G A T A C C T A A C G T G T T A G T
```

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highest scoring alignments shared by a certain number of genomes (using a quorum)

highest scoring alignments supported by a subset of genomes (using a search color set)

# **Alignment Statistics**

Alignment statistic for BLAST:

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BLAST ranks findings according to significance values

E-value:

$$E = \underbrace{Kmn}_{C} e^{-\lambda S}$$
 **P-value**:  
 $P = 1 - e^{-E}$ 

where n = query length, m = database size

## **Alignment Statistics**

Alignment statistic for BLAST:

BLAST ranks findings according to significance values

E-value:

$$E = \underbrace{Kmn}_{C} e^{-\lambda S}$$
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where n = query length, m = database size

How does a sequence to graph alignment statistic work?

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- Graph sequences are highly similar (random hits –)
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- Maximum scores follows an exponential distribution
- =  $\lambda$  and C depend on sequence relatedness and diversity within pangenome

But: Details unknown!

## Parameter estimations

How to get good estimates for  $\lambda$  and C?

Our approach: Obtain estimates by random sampling

Problem:

- Frequency of high scoring alignments is very low ( $< 10^{-6}$ )
- A lot of sampling required
- But: Parameters need to be reliable esp. for high scores

 $\rightarrow$  Importance sampling based on *Metropolis-Hastings Markov Chain Monte Carlo* (MCMC) strategy

## Importance sampling via MCMC

Idea: Construct Markov chain such that the probability to sample a random sequence with score *s* is exponentially biased towards higher scores

Procedure:

- From sequence x propose a new sequence y by substituting/ deleting\*/ inserting\* a single base in x
- **2** Accept proposal with probability  $\min\{1, \exp(\lambda_0 \cdot (s_y s_x))\}$
- If y is accepted make a new proposal for y (for x otherwise)
- ightarrow after 2*n*/3 accepts, procedure yields new sample with uncorrelated score
- $\rightarrow$  Calculate sample's score

 $<sup>^{*}</sup>$  with shifting in/out a base at the sequence borders

# Estimation of $\lambda$ and C

Follow "naive" random sampling strategy to obtain initial estimates for  $\lambda$  and C

Generate many sample sequences using MCMC and score them

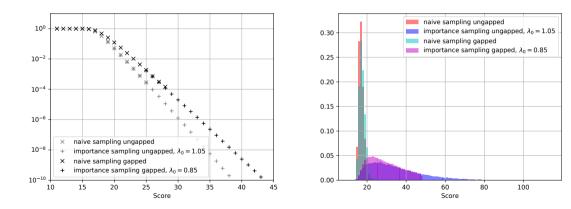
Let  $R_s$  be the absolute number of times score s was observed among all samples, then

$$\mathcal{T}_{s} := \Sigma_{s' \geq s} R_{s'} \cdot \exp(-\lambda_0 \cdot s')$$

 $\lambda$  can be estimated by fitting line to points  $(s, \log T_s)$ 

C is estimated from 10% of highest scores of naive sampling

## Estimation of $\lambda$ and C for 220 Salmonella enterica genomes



18

# Conceptual comparison: BLAST vs. PLAST

### BLAST

- Allows comparison of
  - query and each DB sequence

- Processes every DB sequence independent from each other
- Increased runtime if sequence similarity increases

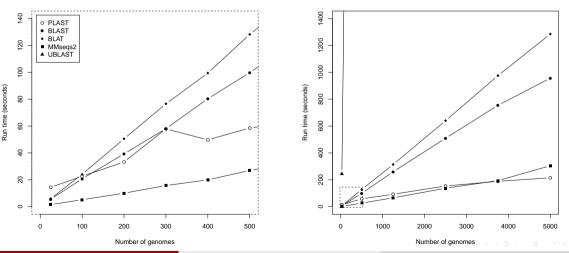
## PLAST

- Allows comparison of
  - query and each graph sequence
  - graph sequences among each other (with respect to query)
- Processes many sequences in parallel most of the time
- Saves run time with increasing similarity of graph sequences

## Runtime and memory comparison

**Database:** 5,000 *Salmonella typhimurium* assemblies (total: 24 GB)

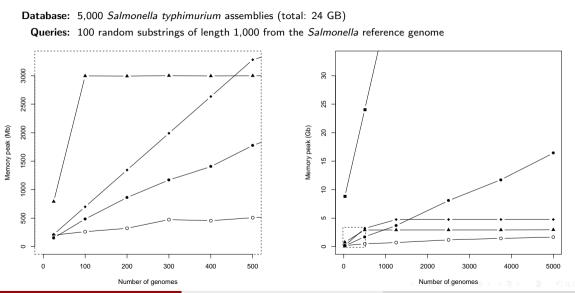
Queries: 100 random substrings of length 1,000 from the Salmonella reference genome



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## Runtime and memory comparison



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## Experimental comparison to other tools

Tool	Results	Tool\PLAST	PLAST\Tool	Tool\BLAST	BLAST\Tool
PLAST	7,565	-	-	357 4.72 %	290 0.02 %
BLAST	1,246,221	290 0.02 %	357 4.72 %	-	-
BLAT	457,089	0.00%	456 6.03 %	$508 \\ 0.11 \%$	49,798 4.00 %
MMseqs2	695,792	6 0.00 %	322 4.26 %	800 0.12 %	$21,022 \\ 1.69\%$
UBLAST	4,881,509	$^{111,386}_{2.28\%}$	272 3.59 %	220,577 4.52 %	5,459 0.44 %

Two results match if they overlap by at least 90 % of the shorter alignment.

#### Not included:

- DIAMOND (protein databases only)
- SWORD (proteins only)
- GHOSTZ (index building took too long)

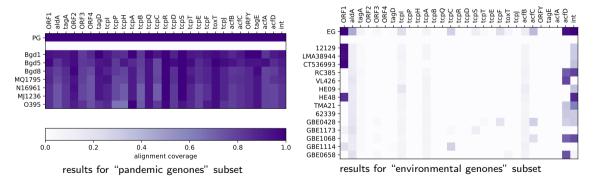
- LAST (index building took too long)
- RAPSearch2 (reported issue)
- LASTZ, YASS, BLASTZ (performed worse in earlier studies)

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Use case: Pathogenicity islands in Vibrio cholerae

**Database:** *Vibrio cholerae* pangenome consisting of 21 genomes; some assembled, some not **Query:** VPI-1 region: 19 genes associated with pathogenicity



## Beyond microbial pangenomes

#### Construction of a human pangenome

We used

- variant information from the 1000 Genomes Project phase 3 and
- human reference genome GRCh37

to construct a pangenome of all 2,504 human individuals for

- chromosome 2 and
- chromosome 15

## Beyond microbial pangenomes

#### Exemplary use case: Investigation of polymorphism within human

#### SNP rs1426654 influences skin pigmentation

- Reference allele: light skin color (West Eurasian ancestry)
- Located on exon 3 of gene SLC24A5 (chr. 15)

### PLAST input:

- Pangenome of human chromosome 15
- Reference sequence of exon 3 as query (GRCh37)
- Result:
  - ▶ No quorum: perfect match (reference), single mismatch (variant)
  - > 99% quorum of all European samples: reference allele exclusively

# Beyond microbial pangenomes

### Performance on human pangenome: PLAST vs. MMseqs2

Setup:

- $\triangleright$  Data set: Pangenome of chr. 2 (pprox 8% of comp. human genome)
- Search of 100 1kb queries (randomly drawn from reference)
- PLAST took
  - 317s per query on average
  - ▶ and 24GB of memory
  - on a single core
- MMseqs2 (on subset of only 1,000 chromosomes) took
  - 339s per query on average
  - ▶ and 259 GB of memory
  - on all available 28 cores
- Sizes of input files on disk:
  - ▶ 9.2 GB (PLAST)
  - 2.2 TB (MMseqs2)

## Conclusion and Outlook

#### **Conclusions:**

Increasing amounts of available DNA sequences motivate the usage of pangenomic approaches

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- Searching withing C-DBGs allows comparison of graph sequences in parallel

#### Thank you!

#### Manuscript on bioRxiv: https://doi.org/10.1101/2020.09.03.280958